

# Network Contention Aware HPC job scheduling with Workload Precongnition

Xu Yang<sup>\*#</sup>, John Jenkins<sup>+</sup>, Misbah Mubarak<sup>+</sup>, Robert Ross<sup>+</sup>, Zhiling Lan<sup>\*</sup>

<sup>\*</sup>Illinois Institute of Technology, Chicago, IL, USA

<sup>+</sup>Argonne National Laboratory, Argonne, IL, USA

<sup>#</sup> xyang56@hawk.iit.edu

# Background

- Scale of supercomputers increases

	<i>NO. of Nodes (core/node)</i>	<i>CPU Type (Clock Speed GHz)</i>	<i>Peak Tflop/s (Gflop/s/core)</i>	<i>Network Topolog y</i>	<i>Link bw (GB/s)</i>	<i>Injection (GB/s)</i>
<i><b>Intrepid (2007)</b></i>	<i><b>40960 (04)</b></i>	<i><b>PowerPC450 (0.85)</b></i>	<i><b>557 (3.5)</b></i>	<i><b>3D torus</b></i>	<i><b>0.425</b></i>	<i><b>2.55</b></i>
<i><b>Mira (2012)</b></i>	<i><b>49152 (16)</b></i>	<i><b>PowerPC A2 (1.6)</b></i>	<i><b>10066 (12.8)</b></i>	<i><b>5D torus</b></i>	<i><b>2.000</b></i>	<i><b>20.00</b></i>

- Concurrently running applications

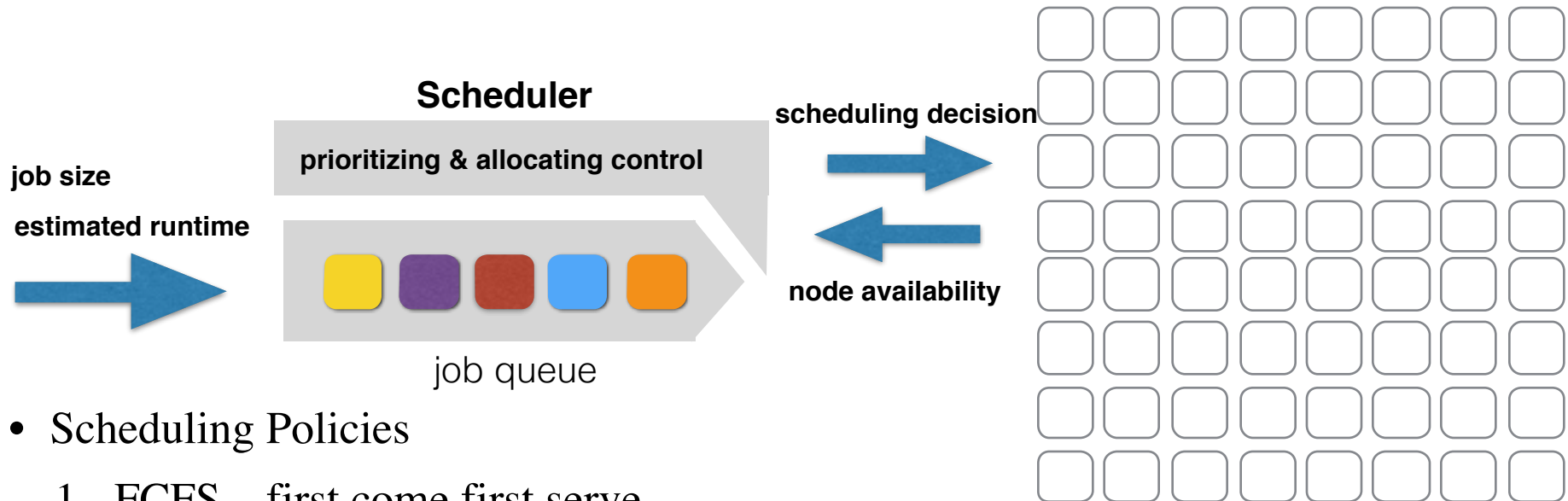
## 1. Network Contention

- allocation shape
- communication pattern

## 2. Performance variability (up to 70%)

# Background

## Traditional batch scheduler



- Scheduling Policies
  1. FCFS—first come first serve
  2. STF—shortest runtime first
  3. FCFS/ Easy backfilling
  4. others based on different priorities

All these policies are priorities based. Priorities could be arrival time, runtime, size, user, etc

# Background

Two types of allocation schemes on Torus system:

1. Non-contiguous (Cray XT)

- no fragmentation
- inter-process communication less efficient
- network contention

2. Partition based scheduling system(IBM BlueGene/P, Q )

- preserve locality
- reduce network contention
- fragmentation and low network utilization

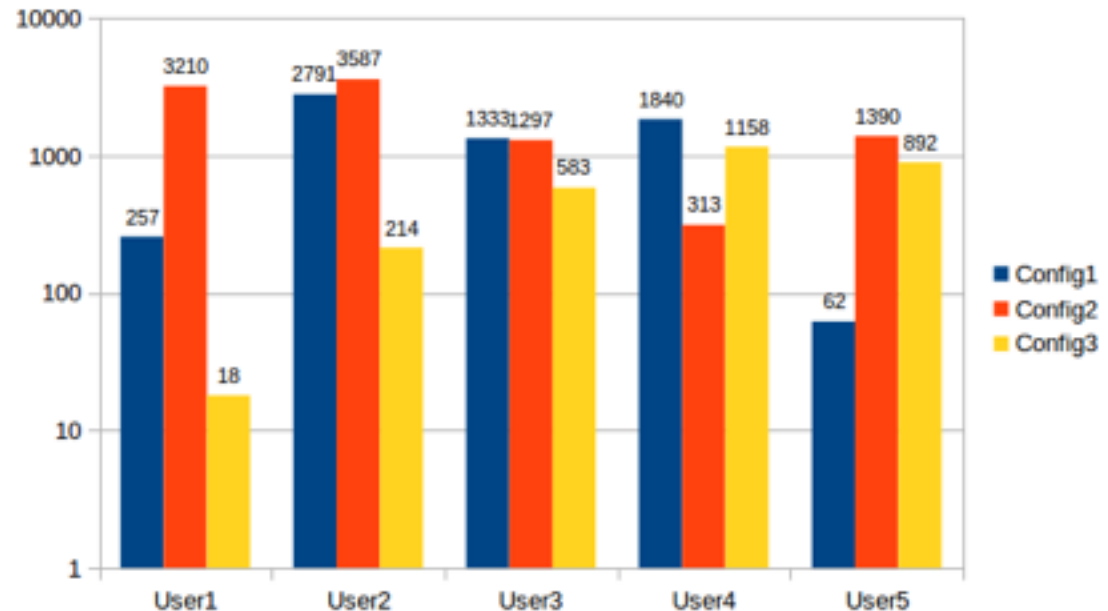
Both have limited knowledge about Job and System:

- job information, such as arrival time, expected runtime, size
- system information, such as network topology, node availability

# Observation

- Job Communication Pattern

1. Nearest Neighbor
2. P2P
3. All to All
4. Collectives



- Repetitive Workload

1. stable user group
2. limited application sets

The repetitive job submissions on Mira in 2014 (March-September). The percentage of repetitive submissions from certain user can be as much as 90%.

# Motivation

- Compact allocation is not necessary for every application, depends on their communication pattern
- Job Scheduler should take jobs' communication information for making scheduling decisions
  - 1.preferable allocation shape
  - 2.better task mapping
- Quantified analysis is needed for
  - 1.bandwidth requirement for job with specific communication pattern
  - 2.network contention between concurrently running jobs
  - 3.performance loss, i.e. extended job run time, lower system throughput

This is where CODES come to rescue!

# CODES Job Allocation

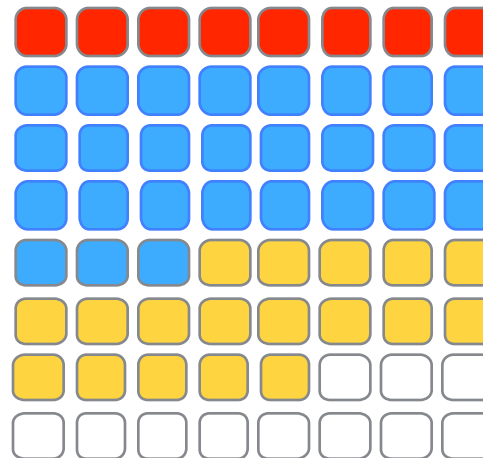
- Things that CODES extension capable of
    1. Concurrently running jobs with given network topology
    2. Locate network contention
    3. Performance variability analysis
  - Workload
    - DOE full apps and associate min-apps
- <http://portal.nersc.gov/project/CAL/designforward.htm>
- Selected Apps
    1. AMG
    2. MiniFE
    3. Crystal Router

# CODES Job Allocation

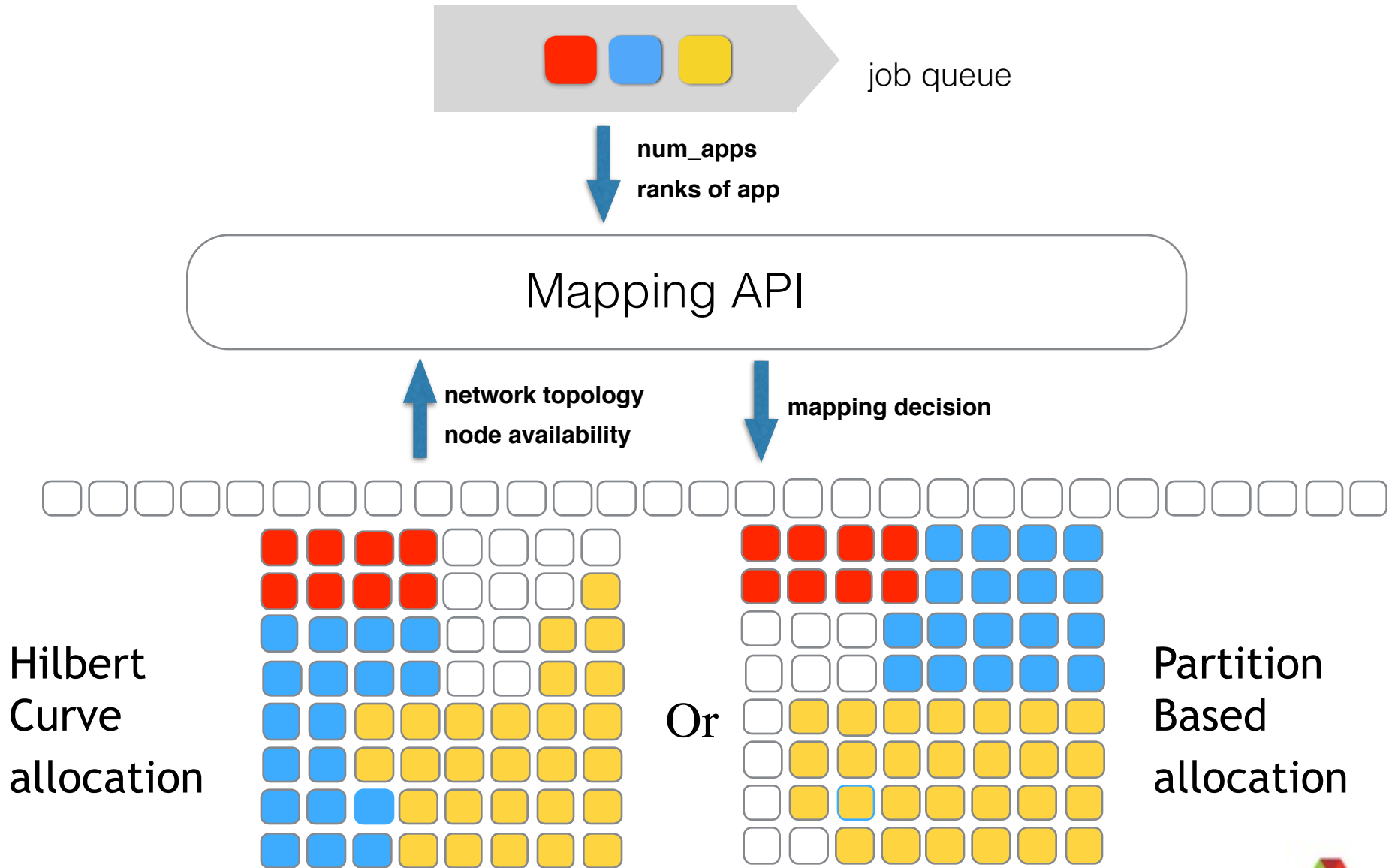
- Network Model  
2D torus



	AMG8	AMG27	MiniFE
NO. of ranks	8	27	18



# CODES Job Allocation



# Communication Pattern V.S Alloc Shape

- All to All

compact allocation with most bandwidth resource

- Nearest Neighbor

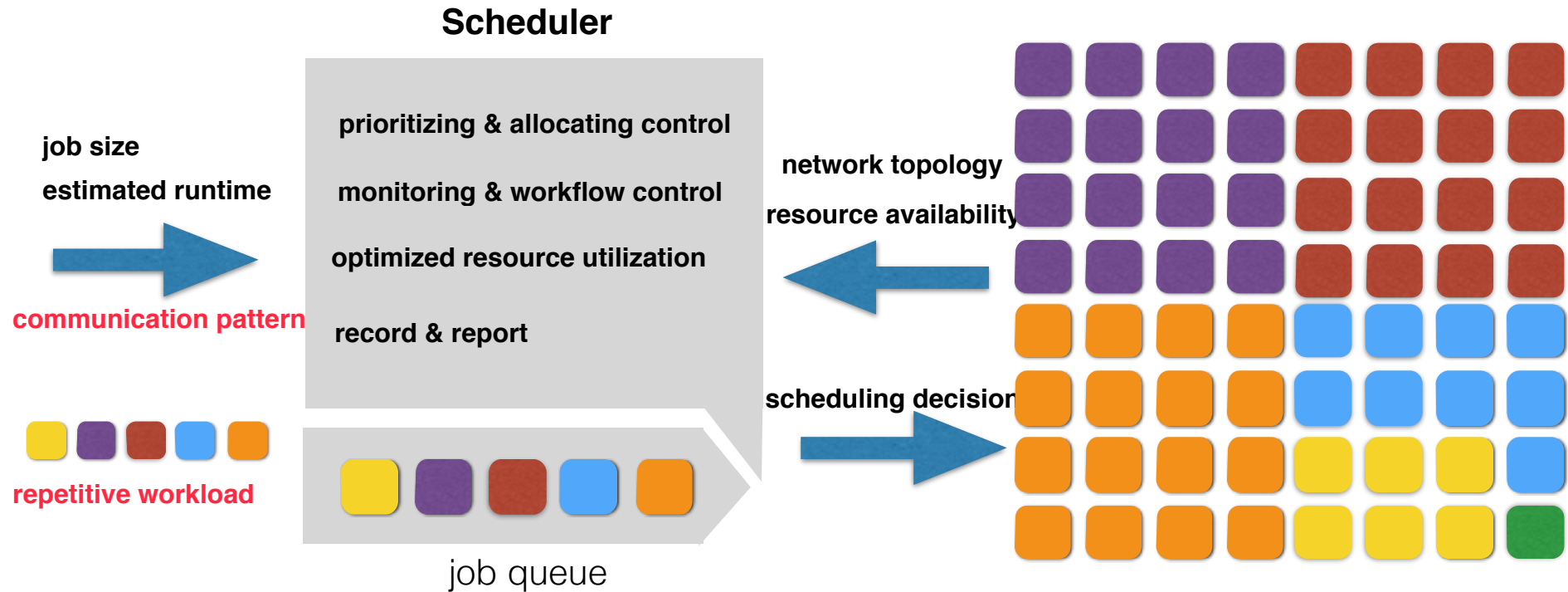
if no compact allocation available, contiguous allocation will be adequate

- P2P and Collectives

without heavy communication, can be flexible

Job with different communication patterns should get different allocation shapes

# New Scheduler Design



# Future Work

- Provide more mapping strategies with different network topologies
- Explore the characteristics of different workloads
- Collective traces from other representative HPC applications